Colouring $(P_r + P_s)$-Free Graphs [*] [†]

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Abstract

The k-Colouring problem is to decide if the vertices of a graph can be coloured with at most k colours for a fixed integer k such that no two adjacent vertices are coloured alike. If each vertex u must be assigned a colour from a prescribed list $L(u) \subseteq \{1, \ldots, k\}$, then we obtain the List k-Colouring problem. A graph G is H-free if G does not contain H as an induced subgraph. We continue an extensive study into the complexity of these two problems for H-free graphs. We prove that List 3-Colouring is polynomial-time solvable for $(P_2 + P_5)$-free graphs and for $(P_3 + P_4)$-free graphs. Combining our results with known results yields complete complexity classifications of 3-Colouring and List 3-Colouring on H-free graphs for all graphs H up to seven vertices. We also prove that 5-Colouring is NP-complete for $(P_3 + P_5)$-free graphs.

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1 Introduction

Graph colouring is a popular concept in Computer Science and Mathematics due to a wide range of practical and theoretical applications, as evidenced by numerous surveys and books on graph colouring and many of its variants (see, for example, [5, 14, 21, 24, 28, 30, 33]). Formally, a *colouring* of a graph $G = (V, E)$ is a mapping $c : V \to \{1, 2, \ldots\}$ that assigns each vertex $u \in V$ a *colour* $c(u)$ in such a way that $c(u) \neq c(v)$ whenever $uv \in E$. If $1 \leq c(u) \leq k$, then c is also called a *k-colouring* of G and G is said to be *k-colourable*. The Colouring problem is to decide if a given graph G has a k-colouring for some given integer k.

It is well known that Colouring is NP-complete even if $k = 3$ [27]. To pinpoint the reason behind the computational hardness of Colouring one may impose restrictions on the input. This led to an extensive study of Colouring for special graph classes, particularly hereditary graph classes. A graph class is *hereditary* if it is closed under vertex deletion. As this is a natural property, hereditary graph classes capture a very large collection of well-studied graph classes. It is readily seen that a graph class $\mathcal{G}$ is hereditary if and only if $\mathcal{G}$ can be characterized by a unique set $\mathcal{H}_{\mathcal{G}}$ of minimal forbidden induced subgraphs. If $\mathcal{H}_{\mathcal{G}} = \{H\}$, then a graph $G \in \mathcal{G}$ is called *H-free*.

Král', Kratochvíl, Tuza, and Woeginger [23] started a systematic study into the complexity of Colouring on $\mathcal{H}$-free graphs for sets $\mathcal{H}$ of size at most 2. They showed polynomial-time solvability if H is an induced subgraph of P_4 or $P_1 + P_3$ and NP-completeness for all other graphs H. The classification for the case where $\mathcal{H}$ has size 2 is far from finished; see the summary in [14] or an updated partial overview in [11] for further details. Instead of considering sets $\mathcal{H}$ of size 2, we consider H-free graphs and follow another well-studied direction, in which the number of colours k is *fixed*, that is, k no longer belongs to the input.

k-Colouring: Given a graph G does there exist a k-colouring of G?

A *k-list assignment* of G is a function L with domain V such that the *list of admissible colours* $L(u)$ of each $u \in V$ is a subset of $\{1, 2, \ldots, k\}$. A colouring c *respects* L if $c(u) \in L(u)$ for every $u \in V$. If k is fixed, then we obtain the following generalization of k-Colouring:

List k-Colouring: Given a graph G and a k-list assignment L does there exist a colouring of G that respects L?

For every $k \geq 3$, k-Colouring on H-free graphs is NP-complete if H contains a cycle [13] or an induced claw [19, 26]. Hence, the case where H is a *linear forest* (a disjoint union of paths) remains. The situation is far from settled yet, although many partial results are known [2, 3, 4, 6, 7, 8, 9, 10, 15, 18, 20, 25, 29, 31, 34]. Particularly, the case where H is the t-vertex path P_t has been well studied. The cases $k = 4$, $t = 7$ and $k = 5$, $t = 6$ are NP-complete [20]. For $k \geq 1$, $t = 5$ [18] and $k = 3$, $t = 7$ [2], even List k-Colouring on P_t-free graphs is polynomial-time solvable (see also [14]). For a fixed integer k, the k-Precolouring Extension problem is to decide a given k-colouring defined on an induced subgraph of a graph G can be extended to a k-colouring of G. Recently it was shown in [7, 8] that 4-Precolouring Extension, and therefore 4-Colouring, is polynomial-time solvable for P_6-free graphs. In contrast, the more general problem List 4-Colouring is NP-complete for P_6-free graphs [15]. See Table 1 for a summary of all these results.

From Table 1 we see that only the cases $k = 3$, $t \geq 8$ are still open, although some partial results are known for k-Colouring for the case $k = 3$, $t = 8$ [9]. The situation when H is a disconnected linear forest $\bigcup P_i$ is less clear. It is known that for every $s \geq 1$, List 3-Colouring is polynomial-time solvable for sP_3-free graphs [4, 14]. For every graph H,

	k-Colouring				k-Precolouring Extension				List k-Colouring			
t	$k=3$	$k=4$	$k=5$	$k\geq 6$	$k=3$	$k=4$	$k=5$	$k\geq 6$	$k=3$	$k=4$	$k=5$	$k\geq 6$
$t\leq 5$	P	P	P	P	P	P	P	P	P	P	P	P
$t=6$	P	P	NP-c	NP-c	P	P	NP-c	NP-c	P	NP-c	NP-c	NP-c
$t=7$	P	NP-c	NP-c	NP-c	P	NP-c	NP-c	NP-c	P	NP-c	NP-c	NP-c
$t\geq 8$	?	NP-c	NP-c	NP-c	?	NP-c	NP-c	NP-c	?	NP-c	NP-c	NP-c

Table 1 Summary for P_t-free graphs.

List 3-Colouring is polynomial-time solvable for $(H+P_1)$-free graphs if it is polynomially solvable for H-free graphs [4, 14]. If $H=rP_1+P_5$ ($r\geq 0$) a stronger result is known.

▶ **Theorem 1** ([10]). *For all $k\geq 1, r\geq 0$,* List k-Colouring *is polynomial-time solvable on (rP_1+P_5)-free graphs.*

Theorem 1 cannot be extended to larger linear forests H, as List 4-Colouring is NP-complete for P_6-free graphs [15] and List 5-Colouring is NP-complete for (P_2+P_4)-free graphs [10]. As mentioned, 5-Colouring is known to be NP-complete for P_6-free graphs [20], but the existence of integers $k\geq 3$ and $2\leq r\leq 5$ such that k-Colouring is NP-complete for (P_r+P_5)-free graphs has not been shown in the literature.

Another way of making progress is to complete a classification by bounding the size of H. It follows from the above results and the ones in Table 1 that for a graph H with $|V(H)|\leq 6$, 3-Colouring and List 3-Colouring (and consequently, 3-Precolouring Extension) are polynomial-time solvable on H-free graphs if H is a linear forest, and NP-complete otherwise; see also [14]. In [14] it was also shown that, to obtain the same statement for graphs H with $|V(H)|\leq 7$, only the two cases where $H\in\{P_2+P_5, P_3+P_4\}$ must be solved.

Our Results In Section 2 we solve the two missing cases listed above.

▶ **Theorem 2.** List 3-Colouring *is polynomial-time solvable for (P_2+P_5)-free graphs and for (P_3+P_4)-free graphs.*

We prove Theorem 2 as follows. If the graph G of an instance (G,L) of List 3-Colouring is P_7-free, then we can use the aforementioned result of Bonomo et al. [2]. Hence we may assume that G contains an induced P_7. We consider every possibility of colouring the vertices of this P_7 and try to reduce each resulting instance to a polynomial number of smaller instances of 2-Satisfiability. As the latter problem can be solved in polynomial time, the total running time of the algorithm will be polynomial. The crucial proof ingredient is that we partition the set of vertices of G that do not belong to the P_7 into subsets of vertices that are of the same distance to the P_7. This leads to several "layers" of G. We analyse how the vertices of each layer are connected to each other and to vertices of adjacent layers so as to use this information in the design of our algorithm.

Combining Theorem 2 with the aforementioned known results yields the following complexity classifications for graphs H up to seven vertices.

▶ **Corollary 3.** *Let H be a graph with $|V(H)|\leq 7$. If H is a linear forest, then* List 3-Colouring *is polynomial-time solvable for H-free graphs; otherwise already* 3-Colouring *is* NP*-complete for H-free graphs.*

In Section 3 we complement Theorem 2 by proving the following result.

▶ **Theorem 4.** 5-Colouring *is* NP*-complete for (P_3+P_5)-free graphs.*

Preliminaries

Let $G = (V, E)$ be a graph. For a vertex $v \in V$, we denote its *neighbourhood* by $N(v) = \{u \mid uv \in E\}$, its *closed neighbourhood* by $N[v] = N(v) \cup \{v\}$ and its degree by $\deg(v) = |N(v)|$. For a set $S \subseteq V$, we write $N(S) = \bigcup_{v \in S} N(v) \setminus S$ and $N[S] = N(S) \cup S$, and we let $G[S] = (S, \{uv \mid u, v \in S\})$ be the subgraph of G induced by S. The *contraction* of an edge $e = uv$ removes u and v from G and introduces a new vertex which is made adjacent to every vertex in $N(u) \cup N(v)$. The *identification* of a set $S \subseteq V$ by a vertex w removes all vertices of S from G, introduces w as a new vertex and makes w adjacent to every vertex in $N(S)$. The *length* of a path is its number of edges. The *distance* $\mathrm{dist}_G(u, v)$ between two vertices u and v is the length of a shortest path between them in G. The *distance* $\mathrm{dist}_G(u, S)$ between a vertex $u \in V$ and a set $S \subseteq V \setminus \{v\}$ is defined as $\min\{\mathrm{dist}(u, v) \mid v \in S\}$.

For two graphs G and H, we use $G + H$ to denote the disjoint union of G and H, and we write rG to denote the disjoint union of r copies of G. Let (G, L) be an instance of List 3-Colouring. For $S \subseteq V(G)$, we write $L(S) = \bigcup_{u \in S} L(u)$. We let P_n and K_n denote the path and complete graph on n vertices, respectively. The *diamond* is the graph obtained from K_4 after removing an edge. We say that an instance (G', L') is *smaller* than some other instance (G, L) of List 3-Colouring if either G' is an induced subgraph of G with $|V(G')| < |V(G)|$; or $G' = G$ and $L'(u) \subseteq L(u)$ for each $u \in V(G)$, such that there exists at least one vertex u^* with $L'(u^*) \subset L(u^*)$.

2 The Two Polynomial-Time Results

In this section we show that List 3-Colouring problem is polynomial-time solvable for $(P_2 + P_5)$-free graphs and for $(P_3 + P_4)$-free graphs. As arguments for these two graph classes are overlapping, we prove both cases simultaneously. Our proof uses the following two results.

▶ **Theorem 5** ([2]). List 3-Colouring *is polynomial-time solvable for P_7-free graphs.*

▶ **Theorem 6** ([12]). *The* 2-List Colouring *problem is linear-time solvable.*

Outline of the proof of Theorem 2. Our goal is to reduce, in polynomial time, an instance (G, L) of List 3-Colouring, where G is $(P_2 + P_5)$-free or $(P_3 + P_4)$-free, to a polynomial number of smaller instances of 2-List-Colouring in such a way that (G, L) is a yes-instance if and only if at least one of the new instances is a yes-instance. As for each of the smaller instances, we can apply Theorem 6, the total running time of our algorithm will be polynomial.

If G is P_7-free, then we do not have to do the above and may apply Theorem 5 instead. Hence, we assume that G contains an induced P_7. We put the vertices of the P_7 in a set N_0 and define sets N_i ($i \geq 1$) of vertices of the same distance i from N_0; we say that the sets N_i are the layers of G. We then analyse the structure of these layers using the fact that G is $(P_2 + P_5)$-free or $(P_3 + P_4)$-free. The first phase of our algorithm is about preprocessing (G, L) after colouring the seven vertices of N_0 and applying a number of propagation rules. We consider every possible colouring of the vertices of N_0. In each branch we may have to deal with vertices u that still have a list $L(u)$ of size 3. We call such vertices active and prove that they all belong to N_2. We then enter the second phase of our algorithm. In this phase we show, via some further branching, that N_1-neighbours of active vertices either all have a list from $\{\{h, i\}, \{h, j\}\}$, where $\{h, i, j\} = \{1, 2, 3\}$, or they all have the same list $\{h, i\}$. In the third phase we reduce, again via some branching, to the situation where only the latter option applies: N_1-neighbours of active vertices all have the same list. Then in the fourth and final phase of our algorithm we know so much structure of the instance that we can reduce to a polynomial number of smaller instances of 2-List-Colouring via a new propagation rule identifying common neighbourhoods of two vertices by a single vertex.

Theorem 2 (restated). List 3-Colouring *is polynomial-time solvable for* $(P_2 + P_5)$*-free graphs and for* $(P_3 + P_4)$*-free graphs.*

Proof Sketch. Due to space limitation we omit the proof for the (more involved) case where $H = P_3 + P_4$. Hence, let (G, L) be an instance of List 3-Colouring, where $G = (V, E)$ is a $(P_2 + P_5)$-free graph. Whenever possible, we base our arguments on $(P_3 + P_5)$-freeness. Since the problem can be solved component-wise, we may assume that G is connected. If G contains a K_4, then G is not 3-colourable, and thus (G, L) is a no-instance. As we can decide if G contains a K_4 in $O(n^4)$ time by brute force, we assume that from now on G is K_4-free. By brute force we either deduce in $O(n^7)$ time that G is P_7-free or we find an induced P_7 on vertices $v_1, \ldots, v_7$ in that order. In the first case we use Theorem 5. It remains to deal with the second case.

Definition (Layers). Let $N_0 = \{v_1, \ldots, v_7\}$. For $i \geq 1$, we define $N_i = \{u \mid \text{dist}(u, N_0) = i\}$. We call the sets N_i $(i \geq 0)$ the *layers* of G.

In the remainder, we consider N_0 to be a fixed set of vertices. That is, we will update (G, L) by applying a number of propagation rules and doing some (polynomial) branching, but we will never delete the vertices of N_0. This will enable us to exploit the H-freeness of G.

We show the following two claims about layers (proofs omitted).

Claim 1. $V = N_0 \cup N_1 \cup N_2 \cup N_3$.

Claim 2. $G[N_2 \cup N_3]$ *is the disjoint union of complete graphs of size at most* 3*, each containing at least one vertex of* N_2 *(and thus at most two vertices of* N_3*).*

We will now introduce a number of propagation rules, which run in polynomial time. We are going to apply these rules on G *exhaustively*, that is, until none of the rules can be applied anymore. Note that during this process some vertices of G may be deleted (due to Rules 4 and 10), but as mentioned we will ensure that we keep the vertices of N_0, while we may update the other sets N_i $(i \geq 1)$. We say that a propagation rule is *safe* if the new instance is a yes-instance of List 3-Colouring if and only if the original instance is so.

Rule 1. **(no empty lists)** If $L(u) = \emptyset$ for some $u \in V$, then return `no`.

Rule 2. **(not only lists of size 2)** If $|L(u)| \leq 2$ for every $u \in V$, then apply Theorem 6.

Rule 3. **(connected graph)** If G is disconnected, then solve List 3-Colouring on each instance (D, L_D), where D is a connected component of G that does not contain N_0 and L_D is the restriction of L to D. If D has no colouring respecting L_D, then return `no`; otherwise remove the vertices of D from G.

Rule 4. **(no coloured vertices)** If $u \notin N_0$, $|L(u)| = 1$ and $L(u) \cap L(v) = \emptyset$ for all $v \in N(u)$, then remove u from G.

Rule 5. **(single colour propagation)** If u and v are adjacent, $|L(u)| = 1$, and $L(u) \subseteq L(v)$, then set $L(v) := L(v) \setminus L(u)$.

Rule 6. **(diamond colour propagation)** If u and v are adjacent and share two common neighbours x and y with $L(x) \neq L(y)$, then set $L(x) := L(x) \cap L(y)$ and $L(y) := L(x) \cap L(y)$.

Rule 7. **(twin colour propagation)** If u and v are non-adjacent, $N(u) \subseteq N(v)$, and $L(v) \subset L(u)$, then set $L(u) := L(v)$.

Rule 8. **(triangle colour propagation)** If u, v, w form a triangle, $|L(u) \cup L(v)| = 2$ and $|L(w)| \geq 2$, then set $L(w) := L(w) \setminus (L(u) \cup L(v))$, so $|L(w)| \leq 1$.

Rule 9. **(no free colours)** If $|L(u) \setminus L(N(u))| \geq 1$ and $|L(u)| \geq 2$ for some $u \in V$, then set $L(u) := \{c\}$ for some $c \in L(u) \setminus L(N(u))$.

Rule 10. **(no small degrees)** If $|L(u)| > |\deg(u)|$ for some $u \in V \setminus N_0$, then remove u from G.

As mentioned, our algorithm will branch at several stages to create a number of new but smaller instances, such that the original instance is a yes-instance if and only if at least one of the new instances is a yes-instance. Unless we explicitly state otherwise, we *implicitly* assume that Rules 1–10 are applied exhaustively immediately after we branch (see also Claim 3). If we apply Rule 1 or 2 on a new instance, then a no-answer means that we will discard the branch. So our algorithm will only return a no-answer for the original instance (G, L) if we discarded all branches. On the other hand, if we can apply Rule 2 on some new instance and obtain a yes-answer, then we can extend the obtained colouring to a colouring of G that respects L, simply by restoring all the already coloured vertices that were removed from the graph due to the rules. We will now state (without proof) Claim 3.

Claim 3. *Rules 1–10 are safe and their exhaustive application takes polynomial time. Moreover, if we have not obtained a yes- or no-answer, then afterwards G is a connected (H, K_4)-free graph, such that $V = N_0 \cup N_1 \cup N_2 \cup N_3$ and $2 \leq |L(u)| \leq 3$ for every $u \in V \setminus N_0$.*

Phase 1. Preprocessing $(\mathbf{G}, \mathbf{L})$

In Phase 1 we will preprocess (G, L) using the above propagation rules. To start off the preprocessing we will branch via colouring the vertices of N_0 in every possible way. By colouring a vertex u, we mean reducing the list of permissible colours to size exactly one. (When $L(u) = \{c\}$, we consider vertex coloured by colour c.) Thus, when we colour some vertex u, we always give u a colour from its list $L(u)$, moreover, when we colour more than one vertex we will always assign distinct colours to adjacent vertices.

Branching I ($O(1)$ branches)

We now consider all possible combinations of colours that can be assigned to the vertices in N_0. That is, we branch into at most 3^7 cases, in which $v_1, \ldots, v_7$ each received a colour from their list. We note that each branch leads to a smaller instance and that (G, L) is a yes-instance if and only if at least one of the new instances is a yes-instance. Hence, if we applied Rule 1 in some branch, then we discard the branch. If we applied Rule 2 and obtained a no-answer, then we discard the branch as well. If we obtained a yes-answer, then we are done. Otherwise we continue by considering each remaining branch separately. For each remaining branch, we denote the resulting smaller instance by (G, L) again.

We will now introduce a new rule, namely Rule 11. We apply Rule 11 together with the other rules. That is, we now apply Rules 1–11 exhaustively. However, each time we apply Rule 11 we first ensure that Rules 1–10 have been applied exhaustively.

Rule 11 **($\mathbf{N_3}$-reduction)** If u and v are in N_3 and are adjacent, then remove u and v from G.

We state (without proofs) the following claims.

Claim 4. *Rule 11, applied after exhaustive application of Rules 1–10, is safe and takes polynomial time. Moreover, afterwards G is a connected (H, K_4)-free graph, such that $V = N_0 \cup N_1 \cup N_2 \cup N_3$ and $2 \leq |L(u)| \leq 3$ for every $u \in V \setminus N_0$.*

Claim 5. *The set N_3 is independent, and moreover, each vertex $u \in N_3$ has $|L(u)| = 2$ and exactly two neighbours in N_2 which are adjacent.*

The following claim follows immediately from Claims 2 and 5.

Claim 6. *Every connected component D of $G[N_2 \cup N_3]$ is a complete graph with either $|D| \leq 2$ and $D \subseteq N_2$, or $|D| = 3$ and $|D \cap N_3| \leq 1$.*

The following claim (proof omitted) describes the location of the vertices with a list of size 3.

Claim 7. *For every $u \in V$, if $|L(u)| = 3$, then $u \in N_2$.*

We will now show how to branch in order to reduce the lists of the vertices $u \in N_2$ with $|L(u)| = 3$ by at least one colour. We formalize this approach in the following definition.

Definition (Active vertices). A vertex $u \in N_2$ and its neighbours in N_1 are called *active* if $|L(u)| = 3$. Let A be the set of all active vertices. Let $A_1 = A \cap N_1$ and $A_2 = A \cap N_2$. We *deactivate* a vertex $u \in A_2$ if we reduce the list $L(u)$ by at least one colour. We *deactivate* a vertex $w \in A_1$ by deactivating all its neighbours in A_2.

Note that every vertex $w \in A_1$ has $|L(w)| = 2$ by Rule 5 applied on the vertices of N_0. Hence, if we reduce $L(w)$ by one colour, all neighbours of w in A_2 become deactivated by Rule 5, and w is removed by Rule 4. For $1 \leq i \leq j \leq 7$, we let $A(i,j) \subseteq A_1$ be the set of active neighbours of v_i that are not adjacent to v_j and similarly, we let $A(j,i) \subseteq A_1$ be the set of active neighbours of v_j that are not adjacent to v_i.

Phase 2. Reduce the number of distinct sets $\mathbf{A(i,j)}$

We will now branch into $O(n^{45})$ smaller instances such that (G, L) is a yes-instance of List 3-Colouring if and only if at least one of these new instances is a yes-instance. Each new instance will have the following property:

(P) for $1 \leq i \leq j \leq 7$ with $j - i \geq 2$, either $A(i,j) = \emptyset$ or $A(j,i) = \emptyset$.

Branching II ($O\big(n^{\left(3\cdot(\binom{7}{2}-6)\right)}\big) = O(n^{45})$ branches)
Consider two vertices v_i and v_j with $1 \leq i \leq j \leq 7$ and $j - i \geq 2$. Assume without loss of generality that v_i is coloured 3 and that v_j is coloured either 1 or 3. Hence, every $w \in A(i,j)$ has $L(w) = \{1,2\}$, whereas every $w \in A(j,i)$ has $L(w) = \{2,q\}$ for $q \in \{1,3\}$. We branch as follows. We consider all possibilities where at most one vertex of $A(i,j)$ receives colour 2 (and all other vertices of $A(i,j)$ receive colour 1) and all possibilities where we choose two vertices from $A(i,j)$ to receive colour 2. This leads to $O(n) + O(n^2) = O(n^2)$ branches. In the branches where at most one vertex of $A(i,j)$ receives colour 2, every vertex of $A(i,j)$ will be deactivated. So Property **(P)** is satisfied for i and j.

Now consider the branches where two vertices x_1, x_2 of $A(i,j)$ both received colour 2. We update $A(j,i)$ accordingly. In particular, afterwards no vertex in $A(j,i)$ is adjacent to x_1 or x_2, as 2 is a colour in the list of each vertex of $A(j,i)$. We now do some further branching for those branches where $A(j,i) \neq \emptyset$. We consider the possibility where each vertex of $N(A(j,i)) \cap A_2$ is given the colour of v_j and all possibilities where we choose one vertex in $N(A(j,i)) \cap A_2$ to receive a colour different from the colour of v_j (we consider both options to colour such a vertex). This leads to $O(n)$ branches. In the first branch, every vertex of $A(j,i)$ will be deactivated. So Property **(P)** is satisfied for i and j.

Now consider a branch where a vertex $u \in N(A(j,i)) \cap A_2$ receives a colour different from the colour of v_j. We will show that also in this case every vertex of $A(j,i)$ will be deactivated. For contradiction, assume that $A(j,i)$ contains a vertex w that is not deactivated after colouring u. As u was in $N(A(j,i)) \cap A_2$, we find that u had a neighbour $w' \in A(j,i)$. As u is coloured with a colour different from the colour of v_j, the size of $L(w')$ is reduced by one (due to Rule 4). Hence w' got deactivated after colouring u, and thus $w' \neq w$. As w is still active, w has a neighbour $u' \in A_2$. As u' and w are still active, u' and w are not adjacent to w' or u. Hence, u, w', v_j, w, u' induce a P_5 in G. As x_1 and x_2 both received colour 2, we find that x_1 and x_2 are not adjacent to each other. Hence, x_1, v_i, x_2 induce a P_3 in G. Recall that all vertices of $A(j,i)$, so also w and w', are not adjacent to x_1 or x_2. As u and u' were still active after colouring x_1 and x_2, we find that u and u' are not adjacent to x_1 or x_2 either. By definition of $A(j,i)$, w and w' are not adjacent to v_i. By definition of $A(i,j)$, x_1 and x_2 are not adjacent to v_j. Moreover, v_i and v_j are non-adjacent, as $j - i \geq 2$. We conclude that G contains an induced $P_3 + P_5$, namely with vertex set $\{x_1, v_i, x_2\} \cup \{u, w', v_j, w, u'\}$, a contradiction. Hence, every vertex of $A(j,i)$ is deactivated. So Property **(P)** is satisfied for i and j also for these branches.

Finally by recursive application of the above procedure for all pairs v_i, v_j such that $1 \leq i \leq j \leq 7$ and $j - i \geq 2$ we get a graph satisfying Property **(P)**.

We now consider each resulting instance from Branching II. We denote such an instance by (G, L) again. Note that vertices from N_2 may now belong to N_3, as their neighbours in N_1 may have been removed due to the branching. The exhaustive application of Rules 1– 11 preserves **(P)** (where we apply Rule 11 only after applying Rules 1–10 exhaustively). Hence (G, L) satisfies **(P)**.

We observe that if two vertices in A_1 have a different list, then they must be adjacent to different vertices of N_0. Hence, by Property **(P)**, at most two lists of $\{\{1,2\},\{1,3\},\{2,3\}\}$ can occur as lists of vertices of A_1. Without loss of generality this leads to two cases: either every vertex of A_1 has list $\{1,2\}$ or $\{1,3\}$ and both lists occur on A_1; or every vertex of A_1 has list $\{1,2\}$ only. In the next phase of our algorithm we reduce, via some further branching, every instance of the first case to a polynomial number of smaller instances of the second case.

Phase 3. Reduce to the case where vertices of $\mathbf{A_1}$ have the same list

Recall that we assume that every vertex of A_1 has list $\{1,2\}$ or $\{1,3\}$. In this phase we deal with the case when both types of lists occur in A_1. We first show, without proof, the following two claims.

Claim 8. *Let $i \in \{1,3,5,7\}$. Then every vertex from $A_1 \cap N(v_i)$ is adjacent to some vertex v_j with $j \notin \{i-1, i, i+1\}$.*

Claim 9. *It holds that $N(A_1) \cap N_0 = \{v_{i-1}, v_i, v_{i+1}\}$ for some $2 \leq i \leq 6$. Moreover, we may assume without loss of generality that v_{i-1} and v_{i+1} have colour 3 and both are adjacent to all vertices of A_1 with list $\{1,2\}$, whereas v_i has colour 2 and is adjacent to all vertices of A_1 with list $\{1,3\}$.*

By Claim 9, we can partition the set A_1 into two (non-empty) sets $X_{1,2}$ and $X_{1,3}$, where $X_{1,2}$ is the set of vertices in A_1 with list $\{1,2\}$ whose only neighbours in N_0 are v_{i-1} and v_{i+1} (which both have colour 3) and $X_{1,3}$ is the set of vertices in A_1 with list $\{1,3\}$ whose only neighbour in N_0 is v_i (which has colour 2).

Our goal is to show that we can branch into at most $O(n^2)$ smaller instances, in which either $X_{1,2} = \emptyset$ or $X_{1,3} = \emptyset$, such that (G, L) is a yes-instance of List 3-Colouring if and only if at least one of these smaller instances is a yes-instance. Then afterwards it suffices to show how to deal with the case where all vertices in A_1 have the same list in polynomial time; this will be done in Phase 4 of the algorithm. We start with the following $O(n)$ branching procedure (in each of the branches we may do some further $O(n)$ branching later on).

Branching III ($O(n)$ branches)

We branch by considering the possibility of giving each vertex in $X_{1,2}$ colour 2 and all possibilities of choosing a vertex in $X_{1,2}$ and giving it colour 1. This leads to $O(n)$ branches. In the first branch we obtain $X_{1,2} = \emptyset$. Hence we can start Phase 4 for this branch. We now consider every branch in which $X_{1,2}$ and $X_{1,3}$ are both nonempty. For each such branch we will create $O(n)$ smaller instances of List 3-Colouring, where $X_{1,3} = \emptyset$, such that (G, L) is a yes-instance of List 3-Colouring if and only if at least one of the new instances is a yes-instance.

Let $w \in X_{1,2}$ be the vertex that was given colour 1 in such a branch. Although by Rule 4 vertex w will need to be removed from G, we make an exception by temporarily keeping w after we coloured it. The reason is that the presence of w will be helpful for analysing the structure of (G, L) after Rules 1–11 have been applied exhaustively (where we apply Rule 11 only after applying Rules 1–10 exhaustively). In order to do this, we first show the following three claims (proofs omitted).

Claim 10. *Vertex w is not adjacent to any vertex in $A_2 \cup X_{1,2} \cup X_{1,3}$.*

Claim 11. *The graph $G[X_{1,3} \cup (N(X_{1,3}) \cap A_2) \cup N_3]$ is the disjoint union of one or more complete graphs, each of which consists of either one vertex of $X_{1,3}$ and at most two vertices of A_2, or one vertex of N_3.*

Claim 12. *For every pair of adjacent vertices s, t with $s \in A_2$ and $t \in N_2$, either t is adjacent to w, or $N(s) \cap X_{1,3} \subseteq N(t)$.*

We now continue as follows. Recall that $X_{1,3} \neq \emptyset$. Hence there exists a vertex $s \in A_2$ that has a neighbour $r \in X_{1,3}$. As $s \in A_2$, we have that $|L(s)| = 3$. Then, by Rule 10, we find that s has at least two neighbours t and t' not equal to r. By Claim 11, we find that neither t nor t' belongs to $X_{1,3} \cup N_3$. We are going to fix an induced 3-vertex path P^s of G, over which we will branch, in the following way.

If t and t' are not adjacent, then we let P^s be the induced path in G with vertices t, s, t' in that order. Suppose that t and t' are adjacent. As G is K_4-free and s is adjacent to r, t, t', at least one of t, t' is not adjacent to r. We may assume without loss of generality that t is not adjacent to r.

First assume that $t \in N_2$. Recall that s has a neighbour in $X_{1,3}$, namely r, and that r is not adjacent to t. We then find that t must be adjacent to w by Claim 12. As $s \in A_2$, we find that s is not adjacent to w by Claim 10. In this case we let P^s be the induced path in G with vertices s, t, w in that order.

Now assume that $t \notin N_2$. Recall that $t \notin N_3$. Hence, t must be in N_1. Then, as $t \notin X_{1,3}$ but t is adjacent to a vertex in A_2, namely s, we find that $t \in X_{1,2}$. Recall that $t' \notin X_{1,3}$. If $t' \in N_1$ then the fact that $t' \notin X_{1,3}$, combined with the fact that t' is adjacent to $s \in A_2$, implies that $t' \in X_{1,2}$. However, by Rule 8 applied on s, t, t', vertex s would have a list of size 1 instead of size 3, a contradiction. Hence, $t' \notin N_1$. As $t' \notin N_3$, this means that $t' \in N_2$. If t' is adjacent to r, then $t \in X_{1,2}$ with $L(t) = \{1, 2\}$ and $r \in X_{1,3}$ with $L(r) = \{1, 3\}$ would have the same lists by Rule 6 applied on r, s, t, t', a contradiction. Hence t' is not adjacent to r. Then, by Claim 12, we find that t' must be adjacent to w. Note that s is not adjacent to w due to Claim 10. In this case we let P^s be the induced path in G with vertices s, t', w in that order. We conclude that either $P^s = tst'$ or $P^s = stw$ or $P^s = st'w$. We are now ready to apply two more rounds of branching.

Branching IV ($O(n)$ branches)
We branch by considering the possibility of removing colour 2 from the list of each vertex in $N(X_{1,3}) \cap A_2$ and all possibilities of choosing a vertex in $N(X_{1,3}) \cap A_2$ and giving it colour 2. In the branch where we removed colour 2 from the list of every vertex in $N(X_{1,3}) \cap A_2$, we obtain that $X_{1,3} = \emptyset$. Hence for that branch we can enter Phase 4. Now consider a branch where we gave some vertex $s \in N(X_{1,3}) \cap A_2$ colour 2. Let $P^s = tst'$ or $P^s = stw$ or $P^s = st'w$. We do some further branching by considering all possibilities of colouring the vertices of P^s that are not equal to the already coloured vertices s and w (should w be a

vertex of P^s) and all possibilities of giving a colour to the vertex from $N(s) \cap X_{1,3}$ (recall that by Claim 11, $|N(s) \cap X_{1,3}| = 1$). This leads to a total of $O(n)$ branches. We claim that in both branches, $|X_{1,3}|$ has reduced to at most 1 (proof omitted).

Branching V ($O(1)$ branches)
We branch by considering both possibilities of colouring the unique vertex of $X_{1,3}$. This leads to two new but smaller instances of List 3-Colouring, in each of which the set $X_{1,3} = \emptyset$. Hence, our algorithm can enter Phase 4.

Phase 4. Reduce to a set of instances of 2-List Colouring
Recall that in this stage of our algorithm we have an instance (G, L) in which every vertex of A_1 has the same list, say $\{1, 2\}$. As G is $(P_2 + P_5)$-free, $G[N_2 \cup N_3]$ is an independent set; otherwise two adjacent vertices of $N_2 \cup N_3$ form, together with $v_1, \ldots, v_5$, an induced $P_2 + P_5$. Hence, we can safely colour each vertex in A_2 with colour 3, and afterwards we may apply Theorem 6.

The correctness of our algorithm follows from the description. The branching in the five stages (Branching I-V), yields a total number of $O(n^{47})$ branches and each branch we created takes polynomial time to process. Hence, the running time of our algorithm is polynomial. ◀

Remark. Except for Phase 4 of our algorithm, all arguments in our proof hold for $(P_3 + P_5)$-free graphs. The difficulty in Phase 4 is that in contrary to the previous phases we cannot use the vertices from N_0 to find an induced $P_3 + P_5$ and therefore obtain the contradiction.

3 The Hardness Result

We show that 5-Colouring is NP-complete for $(P_3 + P_5)$-free graphs by reducing from the NP-complete problem [32] Not-All-Equal 3-Satisfiability with positive literals only, defined as follows: given a set $X = \{x_1, x_2, ..., x_n\}$ of logical variables and a set $\mathcal{C} = \{C_1, C_2, ..., C_m\}$ of 3-literal clauses over X in which all literals are positive, is there a truth assignment for X such that each clause contains at least one true literal and at least one false literal? We call such a truth assignment *satisfying*.

Theorem 4 (restated). 5-Colouring *is* NP*-complete for* $(P_3 + P_5)$*-free graphs.*

Proof Sketch. From a given instance $(\mathcal{C}, X)$ of Not-All-Equal 3-Satisfiability with positive literals only, we first construct a graph G with a list assignment L. For each $x_i \in X$ we introduce two vertices x_i and $\overline{x_i}$, which we make adjacent to each other. We say that x_i and $\overline{x}_i$ are of x-type. We set $L(x_i) = L(\overline{x}_i) = \{4, 5\}$. For each $C_j \in \mathcal{C}$ we introduce a vertex C_j and a vertex C'_j called the *copy* of C_j. We say that C_j and C'_j are of C-type. We set $L(C_j) = L(C'_j) = \{1, 2, 3\}$. We add an edge between each x-type vertex and each C-type vertex. For each $C_j \in \mathcal{C}$ we do as follows. We fix an arbitrary order of the literals in C_j. Say $C_j = \{x_g, x_h, x_i\}$ in that order. Then we add six vertices $a_{g,j}$, $a_{h,j}$, $a_{i,j}$, $a'_{g,j}$, $a'_{h,j}$, $a'_{i,j}$ and edges $x_g a_{g,j}$, $a_{g,j} C_j$, $x_h a_{h,j}$, $a_{h,j} C_j$, $x_i a_{i,j}$, $a_{i,j} C_j$ and also edges $\overline{x}_g a'_{g,j}$, $a'_{g,j} C'_j$, $\overline{x}_h a'_{h,j}$, $a'_{h,j} C'_j$, $\overline{x}_i a'_{i,j}$, $a'_{i,j} C'_j$. We say that $a_{g,j}$, $a_{h,j}$, $a_{i,j}$, $a'_{g,j}$, $a'_{h,j}$, $a'_{i,j}$ are of a-type. We set $L(a_{g,j}) = L(a'_{g,j}) = \{1, 4\}$, $L(a_{h,j}) = L(a'_{h,j}) = \{2, 4\}$ and $L(a_{i,j}) = L(a'_{i,j}) = \{3, 4\}$.

We now extend G into a graph G' by adding a clique consisting of five new vertices $k_1, \ldots, k_5$, which we say are of k-type, and by adding an edge between a vertex k_ℓ and a vertex $u \in V(G)$ if and only if $\ell \notin L(u)$. We can show that $(\mathcal{C}, X)$ has a satisfying truth assignment if and only if G' has a 5-colouring, and moreover that G' is $(P_3 + P_5)$-free (proof omitted). As 5-Colouring belongs to NP, this proves the theorem. ◀

4 Conclusions

By solving two new cases we completed the complexity classifications of 3-Colouring and List 3-Colouring on H-free graphs for graphs H up to seven vertices. We showed that both problems become polynomial-time solvable if H is a linear forest, while they stay NP-complete in all other cases. Recall that k-Colouring ($k \geq 3$) is NP-complete on H-free graphs whenever H is not a linear forest. For the case where H is a linear forest, our new NP-hardness result for 5-Colouring for $(P_3 + P_5)$-free graphs bounds, together with the known NP-hardness results of [20] for 4-Colouring for P_7-free graphs and 5-Colouring for P_6-free graphs, the number of open cases of k-Colouring from above.

For future research we aim to our extend our results. In fact we still do not know if there exists a linear forest H such that 3-Colouring is NP-complete for H-free graphs. This is, however, a notorious open problem studied in many papers; for a recent discussion see [16]. It is also open for List 3-Colouring, where an affirmative answer to one of the two problems yields an affirmative answer to the other one [15]. For $k \geq 4$, we emphasize that all open cases involve linear forests H whose connected components are small. For instance, if H has at most six vertices, then the polynomial-time algorithm for 4-Precolouring Extension on P_6-free graphs [7, 8] implies that there are only three graphs H with $|V(H)| \leq 6$ for which we do not know the complexity of 4 Colouring on H-free graphs, namely $H \in \{P_1 + P_2 + P_3, P_2 + P_4, 2P_3\}$ (see [14]).

The main difficulty to extend the known complexity results is that hereditary graph classes characterized by a forbidden induced linear forest are still not sufficiently well understood due to their rich structure. We need a better understanding of these graph classes to make further progress on a wide range of problems. For example, Independent Set is polynomial-time solvable for P_6-free graphs [17], but it is not known if there exists a linear forest H such that it is NP-complete for H-free graphs. A similar situation holds for Odd Cycle Transversal and Feedback Vertex Set and many other problems; see [1] for a survey.

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